

A Monitoring Tool for a Branching-Time Logic

Duncan Attard¹ Adrian Francalanza¹

¹CS, ICT, University of Malta, Malta
{duncan.attard.01, adrian.francalanza}@um.edu.mt

30th September 2016

Property Specification and Monitor Synthesis

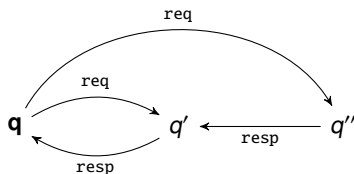
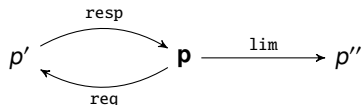
Implementation of a Tool in Erlang

Demo

Modelling Process Behaviour

The behaviour of systems can be described using LTSes that model process execution graphs

Example (Two simple servers)



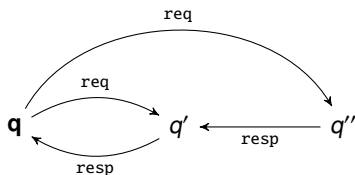
The Branching-Time Logic μHML

Syntax

$$\begin{aligned} \varphi, \phi \in \mu\text{HML} ::= & \mathbf{tt} & | [\alpha]\varphi & | \varphi \wedge \phi & | \mathbf{max} X.\varphi & | X \\ & & | \mathbf{ff} & | \langle \alpha \rangle \varphi & | \varphi \vee \phi & | \mathbf{min} X.\varphi \end{aligned}$$

Example (A liveness property)

$[\text{req}]\langle \text{resp} \rangle \mathbf{tt}$



Cheat sheet

$[\alpha]\mathbf{ff}$ “cannot do event α ”

$\langle \alpha \rangle \mathbf{tt}$ “can do event α ”

A Monitorable Subset

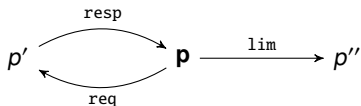
The Syntax

$$\psi \in \text{mHML} \stackrel{\text{def}}{=} \text{sHML} \cup \text{cHML}$$

$$\begin{aligned} \theta, \vartheta \in \text{sHML} &::= \mathbf{tt} \quad | \quad \mathbf{ff} \quad | \quad [\alpha]\theta \quad | \quad \theta \wedge \vartheta \quad | \quad \mathbf{max} X.\theta \quad | \quad X \\ \pi, \varpi \in \text{cHML} &::= \mathbf{tt} \quad | \quad \mathbf{ff} \quad | \quad \langle \alpha \rangle \pi \quad | \quad \pi \vee \varpi \quad | \quad \mathbf{min} X.\pi \quad | \quad X \end{aligned}$$

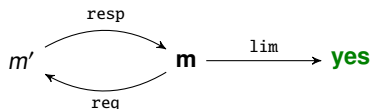
Example (A co-safety property)

$$\mathbf{min} X.(\langle \text{req} \rangle \langle \text{resp} \rangle X \vee \langle \text{lim} \rangle \mathbf{tt})$$



Monitors as LTSes

Example (A monitor for process p)

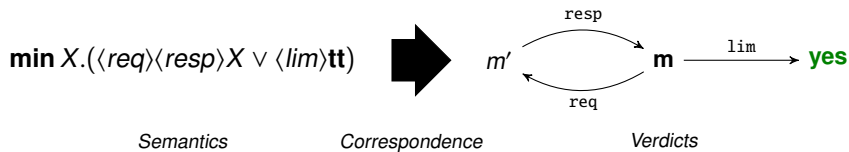


- ▶ A **subtle** difference: processes *generate* actions, monitors *analyse* these actions, and yield verdicts
- ▶ Conclusive verdicts: **no**, **yes**
- ▶ Inconclusive verdict: **end**

Monitor Synthesis

From mHML to Monitors

Correct Synthesis (Compositional)



Monitor Synthesis

From mHML to Monitors

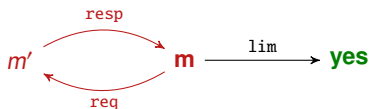
Correct Synthesis (Compositional)

$\min X.(\langle req \rangle \langle resp \rangle X \vee \langle lim \rangle \mathbf{tt})$

Semantics



Correspondence

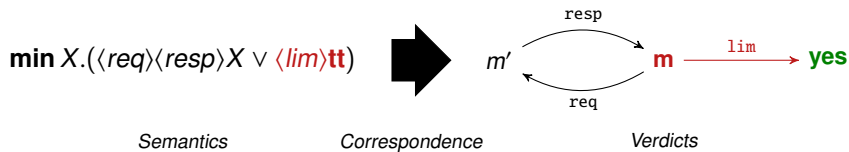


Verdicts

Monitor Synthesis

From mHML to Monitors

Correct Synthesis (Compositional)



Property Specification and Monitor Synthesis

Implementation of a Tool in Erlang

Demo

Actions = Events with Data

- ▶ Why? Because we want to look at the *data* inside actions
- ▶ Events with structure permit us to use pattern matching
- ▶ Output events: $\alpha = \text{client} ! \{\text{resp}, 5\}$
- ▶ Input events: $\alpha = \text{server} ? \{\text{req}, \text{client}, 5\}$

Actions = Events with Data

- ▶ Why? Because we want to look at the *data* inside actions
- ▶ Events with structure permit us to use pattern matching
- ▶ Output events: $\alpha = \text{client} ! \{\text{resp}, 5\}$
- ▶ Input events: $\alpha = \text{server} ? \{\text{req}, \text{client}, 5\}$

Example (Pattern variable binding)

A successor server *Srv* adds one to any numeric payload it receives from clients *Cl*

$$[Srv ? \{\text{req}, Clt, Num\}][Cl ! \{\text{resp}, Num\}]\mathbf{ff}$$

Safety formula ensures that clients *do not* receive the same value *Num* sent by them to the server

The Tool Synthesis

Example (Non-recursive formula)

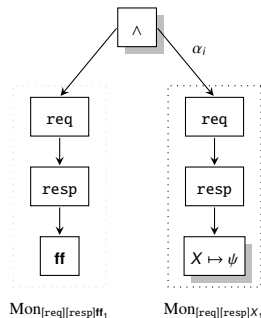
$[Srv ? \{req, Clt, Num\}][Clc ! \{resp, Num\}]\mathbf{ff}$

- ▶ Non-recursive formulae can only observe one interaction before terminating

The Tool Synthesis

Example (Recursive formula)

max $X.($
 $[Srv ? \{req, Clt, Num\}][Cl t ! \{resp, Num\}]ff$
 \wedge
 $[Srv ? \{req, Clt, Num\}][Cl t ! \{resp, Succ\}]X$
 $)$



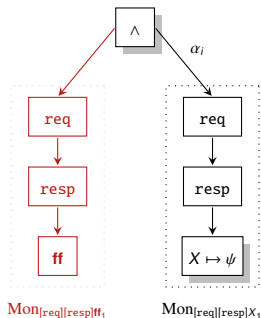
- ▶ Recursive formulae allow continuous monitoring
- ▶ The verdict branch matches events that lead to a detection
- ▶ The recursive branch permits the monitor to unfold *lazily*

The Tool Synthesis

Example (Recursive formula)

$\max X. ($
 $[Srv ? \{req, Clt, Num\}][Cl t ! \{resp, Num\}]ff$
 \wedge
 $[Srv ? \{req, Clt, Num\}][Cl t ! \{resp, Succ\}]X$
 $)$

- ▶ Recursive formulae allow continuous monitoring
- ▶ The **verdict branch** matches events that lead to a detection
- ▶ The recursive branch permits the monitor to unfold *lazily*

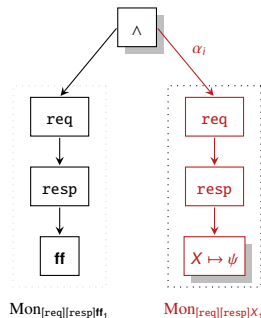


The Tool Synthesis

Example (Recursive formula)

max $X.($
 $[Srv ? \{req, Clt, Num\}][Cl t ! \{resp, Num\}]ff$
 \wedge
 $[Srv ? \{req, Clt, Num\}][Cl t ! \{resp, Succ\}]X$
 $)$

- ▶ Recursive formulae allow continuous monitoring
- ▶ The verdict branch matches events that lead to a detection
- ▶ The **recursive branch** permits the monitor to unfold *lazily*

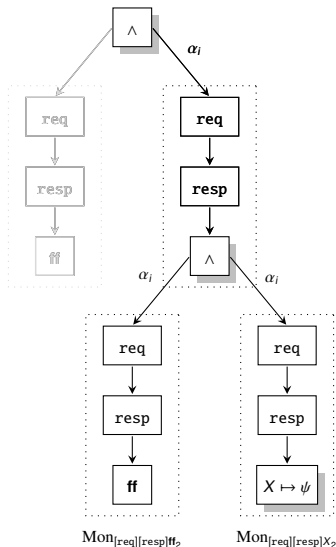


The Tool Synthesis

Example (Recursive formula)

max $X.($
 $[Srv ? \{req, Clt, Num\}][Cl t ! \{resp, Num\}]ff$
 \wedge
 $[Srv ? \{req, Clt, Num\}][Cl t ! \{resp, Succ\}]X$
 $)$

- ▶ Recursive formulae allow continuous monitoring
- ▶ The verdict branch matches events that lead to a detection
- ▶ The recursive branch permits the monitor to unfold *lazily*



Formal vs. Tool Synthesis

$[Srv ? \{req, Clt, Num\}][Cl ! \{resp, Num\}]ff$

```
formula:mon_nec(  
  fun(Action) ->  
    case Action of  
      {recv, Srv, {req, Clt, Num}} ->  
        formula:mon_nec(  
          fun(Action) ->  
            case Action of  
              {send, Clt, {resp, Num}} -> formula:mon_ff();  
              _ -> formula:mon_id()  
            end  
          end);  
      _ -> formula:mon_id()  
    end  
end)
```

Formal vs. Tool Synthesis

$[Srv ? \{req, Clt, Num\}][Clt ! \{resp, Num\}]ff$

```
formula:mon_nec(  
  fun(Action) ->  
    case Action of  
      {recv, Srv, {req, Clt, Num}} ->  
        formula:mon_nec(  
          fun(Action) ->  
            case Action of  
              {send, Clt, {resp, Num}} -> formula:mon_ff();  
              _ -> formula:mon_id()  
            end  
          end);  
      _ -> formula:mon_id()  
    end  
  end)
```

Formal vs. Tool Synthesis

$[Srv ? \{req, Clt, Num\}][Clt ! \{resp, Num\}]ff$

```
formula:mon_nec(  
  fun(Action) ->  
    case Action of  
      {recv, Srv, {req, Clt, Num}} ->  
        formula:mon_nec(  
          fun(Action) ->  
            case Action of  
              {send, Clt, {resp, Num}} -> formula:mon_ff();  
              _ -> formula:mon_id()  
            end  
          end);  
      _ -> formula:mon_id()  
    end  
end)
```

Formal vs. Tool Synthesis

$[Srv ? \{req, Clt, Num\}][Cl ! \{resp, Num\}]ff$

```
formula:mon_nec(  
  fun(Action) ->  
    case Action of  
      {recv, Srv, {req, Clt, Num}} ->  
        formula:mon_nec(  
          fun(Action) ->  
            case Action of  
              {send, Clt, {resp, Num}} -> formula:mon_ff();  
              _ -> formula:mon_id()  
            end  
          end);  
      _ -> formula:mon_id()  
    end  
end)
```

Property Specification and Monitor Synthesis

Implementation of a Tool in Erlang

Demo

Thank You!

Resources

The tool can be downloaded from:

<https://bitbucket.org/duncanatt/detector-lite>

Information and publications from:

<http://www.cs.um.edu.mt/svrg/Tools/detectEr>

Questions?